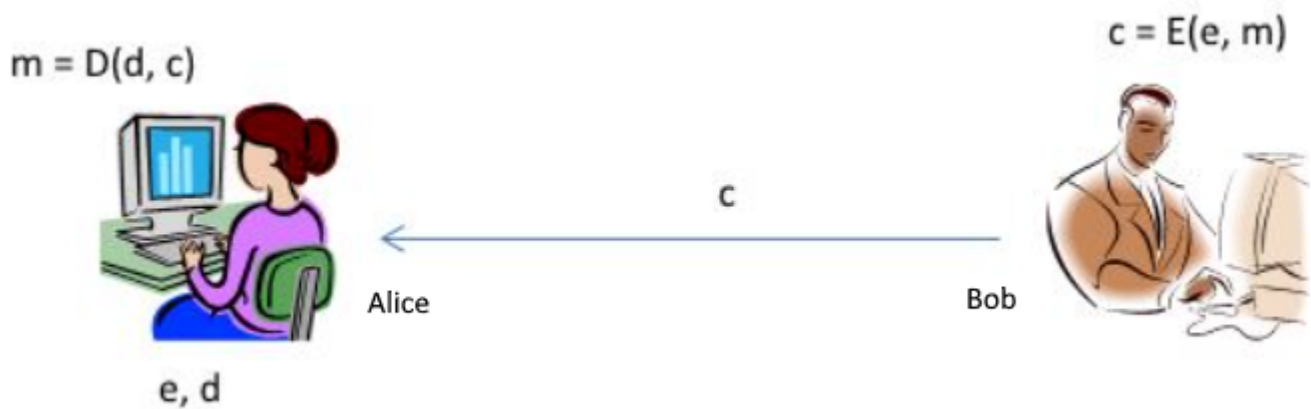


## The Basic Model of Public Key Systems

### Basic communication model



1. Alice generates a pair of keys: **e** (public key) and **d** (private key). In this context **e** means (encryption key) and **d** (decryption key).
2. She keeps **d** secret, but makes **e** public.
3. If Bob wants to send a message to Alice, he uses Alice's public key **e**.
4. Based on the equation  $( c = E(e, m) )$ , only Alice can decrypt  $( c )$  using her private key, with  $( m = D(d, c) )$ , where  $( m )$  is the message.
5. If anyone else wants to send a message to Alice, they can also use her public key **e**.

The system is secure from a decryption perspective because only Alice can decrypt the message, but Alice can never be sure if Bob sent the message, as the public key **e** can be used by anyone.

## RSA

**Rivest, Shamir, Adleman (1977)** developed the RSA algorithm, which is based on exponentiation and modular arithmetic (remainder division). The RSA encryption relies on the fact that if the following equation holds:

$$( T^d \pmod N = T )$$

Then the equation can be split into two parts, where the first equation represents encryption and the second represents decryption:

- **Encryption:**  $( C = T^e \pmod N )$
- **Decryption:**  $( T = C^d \pmod N )$

However, this relationship does not work for just any arbitrary choice of  $( e )$ ,  $( d )$ , and  $( N )$ . The algorithm relies on specific conditions for these values. Let's explore what conditions are required for the RSA algorithm to function correctly.

## How RSA Works

The RSA algorithm uses both **public** and **private keys**. For it to work correctly, the following steps are followed:

### 1.) Key Generation:

- Choose two large prime numbers:  $(p)$  and  $(q)$ .
- Calculate  $(N = p \times q)$ .
- Compute Euler's totient function  $(\phi(N) = (p-1) \times (q-1))$ . ( [totient function](#) )
- Choose a public exponent  $(e)$  such that  $(e)$  is relatively prime to  $(\phi(N))$  (i.e.,  $(1 < e < \phi(N))$  and  $(\gcd(e, \phi(N)) = 1)$ ).
- Compute the private key  $(d)$ , which is the modular multiplicative inverse of  $(e)$ , meaning  $(d \times e \equiv 1 \pmod{\phi(N)})$ .

### 2.) Encryption:

The message  $(T)$  is encrypted using the public key:  $(C = T^e \pmod N)$ .

### 3. Decryption:

The ciphertext  $(C)$  is decrypted using the private key:  $(T = C^d \pmod N)$ .

## Conditions for the Keys

The equations work correctly only when specific conditions are met:

- $(N)$  must be large enough to prevent attackers from easily factoring  $(N)$  into  $(p)$  and  $(q)$ , which would compromise the encryption.
- The public exponent  $(e)$  and the private exponent  $(d)$  must have a specific relationship with  $(N)$  and Euler's totient  $(\phi(N))$ . The relationship between  $(e)$  and  $(d)$  is that  $(e \times d \equiv 1 \pmod{\phi(N)})$ , ensuring that decryption reverses the encryption process.

## Summary

The RSA algorithm depends on the correct choice of prime numbers  $(p)$  and  $(q)$ , the calculation of  $(N)$ , and the mathematical relationship between the exponents  $(e)$  (public key) and  $(d)$  (private key). Without these specific conditions, the encryption and decryption process will not work properly.

Your original description was generally correct, but the conditions governing  $(e)$ ,  $(d)$ , and  $(N)$ , especially the relationship with Euler's totient function, are crucial for the RSA algorithm to function securely.

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